Namn:_____

Person-nummer:_____

Datorarkitektur, 2005

Tentamen 2005-04-02

Instructions:

- Make sure that your exam is not missing any sheets, then write your name and person-nummer on the front. If you need extra pages be sure to write on those too.
- Write your answers in the space provided below the problem. If you make a mess, clearly indicate your final answer.
- The exam has a maximum score of 60 points.
- The problems are of varying difficulty. The point value of each problem is indicated. Pile up the easy points quickly and then come back to the harder problems.
- This exam is OPEN BOOK. You may use any books or notes you like. Good luck!

Problem 1. (10 points):

In the following questions assume the variable a is a signed integer and that the machine uses two's complement representation. Also assume that MAX_INT is the maximum integer, MIN_INT is the minimum integer, and W is one less than the word length (e.g., W = 31 for 32-bit integers). Match each of the descriptions on the left with a line of code on the right (write in the letter). You will be given 2 points for each correct match.

2. Two's complement of a. b. $((a < 0) ? a + 1 : a) >> 1$	1. a%2.	a.	~((a (~a + 1)) >> W) & 1
3. a^a . 4. $!a$. 5. $(a > 0) ? 1 : 0$. 3. a^a . 4. $!((a >> W) !a)$ 6. $!((a^0) << 1)) \& a$ 6. $(((a^0) << 1)) \& a$ 7. $(MAX_INT \& a) + ((a^0))$ 7. $(a > 0) ? 1 : 0$.	2. Two's complement of a.	b.	((a < 0) ? a + 1 : a) >> 1
$d. \ !((a >> W) !a)$ $e. \ (~((~0) << 1)) \& a$ $f. \ (MAX_INT \& a) + (~0)$ $5. \ (a > 0) ? 1 : 0.$		c.	~a - (MIN_INT + MAX_INT)
<pre>4. !a f. (MAX_INT & a) + (~0) 5. (a > 0) ? 1 : 0.</pre>	3. a^a.	d.	!((a >> W) !a)
5. (a > 0) ? 1 : 0.	4. !a.	e.	(~((~0) << 1)) & a
	5. (a > 0) ? 1 : 0.	f.	(MAX_INT & a) + (~0)
		g.	~((a ^ (~0)) a)

Problem 2. (12 points):

Consider the following 16-bit floating point representation based on the IEEE floating point format:

- There is a sign bit in the most significant bit.
- The next seven bits are the exponent. The exponent bias is 63.
- The last eight bits are the significand.

The rules are like those in the IEEE standard (normalized, denormalized, representation of 0, infinity, and NAN).

We consider the floating point format to encode numbers in a form:

$$(-1)^s \times m \times 2^E$$

where m is the *mantissa* and E is the exponent.

Fill in the table below for the following numbers, with the following instructions for each column:

Hex: The 4 hexadecimal digits describing the encoded form.

m: The fractional value of the mantissa. This should be a number of the form x or x/y, where x is an integer, and y is an integral power of 2. Examples include: 0, 67/64, and 1/256.

E: The integer value of the exponent.

Value: The numeric value represented. Use the notation x or $x \times 2^{z}$, where x and z are integers.

As an example, to represent the number 7/2, we would have s = 0, m = 7/4, and E = 1. Our number would therefore have an exponent field of 0×40 (decimal value 63 + 1 = 64) and a significand field $0 \times C0$ (binary 1100000_2), giving a hex representation 40C0.

You need not fill in entries marked "—".

Description	Hex	m	E	Value
-0				-0
Smallest value > 1				
Largest Denormalized				
$-\infty$		_	_	$-\infty$
Number with hex representation 3AA0	3AA0			

Problem 3. (8 points): Consider a **5-bit** two's complement representation. Fill in the empty boxes in the following table. Addition and subtraction should be performed based on the rules for 5-bit, two's complement arithmetic

Number	Decimal Representation	Binary Representation
Zero	0	
n/a	-2	
n/a	9	
n/a	-14	
n/a		0 1100
n/a		1 0100
TMax		
TMin		
TMin+TMin		
TMin+1		
TMax+1		
-TMax		
-TMin		

Problem 4. (8 points):

Consider the source code below, where M and N are constants declared with #define.

```
int mat1[M][N];
int mat2[N][M];
int sum_element(int i, int j)
{
  return mat1[i][j] + mat2[i][j];
}
```

A. Suppose the above code generates the following assembly code:

```
sum_element:
    pushl %ebp
    movl %esp,%ebp
    movl 8(%ebp),%eax
    movl 12(%ebp),%eax
    sall $2,%ecx
    leal 0(,%eax,8),%edx
    subl %eax,%edx
    leal (%eax,%eax,4),%eax
    movl mat2(%ecx,%eax,4),%eax
    addl mat1(%ecx,%edx,4),%eax
    movl %ebp,%esp
    popl %ebp
    ret
```

What are the values of M and N?

```
M =
N =
```

Problem 5. (4 points):

Consider the following C functions and assembly code:

```
int fun4(int *ap, int *bp)
{
    int a = *ap;
    int b = *bp;
    return a+b;
}
                                         pushl %ebp
                                         movl %esp,%ebp
int fun5(int *ap, int *bp)
                                        movl 8(%ebp),%edx
                                        movl 12(%ebp),%eax
{
    int b = *bp;
                                        movl %ebp,%esp
    *bp += *ap;
                                         movl (%edx),%edx
    return b;
                                         addl %edx,(%eax)
}
                                        movl %edx,%eax
                                        popl %ebp
int fun6(int *ap, int *bp)
                                        ret
{
    int a = *ap;
    *bp += *ap;
    return a;
}
Which of the functions compiled into the assembly code shown?
```

Problem 6. (8 points):

Condider the following assembly code for a C for loop:

```
loop:
        pushl %ebp
        movl %esp,%ebp
        movl 8(%ebp),%ecx
        movl 12(%ebp),%edx
        xorl %eax,%eax
        cmpl %edx,%ecx
        jle .L4
.L6:
        decl %ecx
        incl %edx
        incl %eax
        cmpl %edx,%ecx
        jg .L6
.L4:
        incl %eax
        movl %ebp,%esp
        popl %ebp
        ret
```

Based on the assembly code above, fill in the blanks below in its corresponding C source code. (Note: you may only use the symbolic variables x, y, and result in your expressions below — *do not use register names*.)

Problem 7. (10 points):

Consider the following incomplete definition of a C struct along with the incomplete code for a function func given below.

typedef struct node {	node_t n;
x <i>i</i>	<pre>void func() {</pre>
y;	node_t *m;
struct node *next;	m =;
struct node *prev;	m->y /= 16;
} node_t;	return; }

When this C code was compiled on an IA-32 machine running Linux, the following assembly code was generated for function func.

func:

```
pushl %ebp
movl n+12,%eax
movl 16(%eax),%eax
movl %esp,%ebp
movl %ebp,%esp
shrw $4,8(%eax)
popl %ebp
ret
```

Given these code fragments, fill in the blanks in the C code given above. Note that there is a unique answer.

Туре	Size (bytes)	Alignment (bytes)
char	1	1
short	2	2
unsigned short	2	2
int	4	4
unsigned int	4	4
double	8	4

The types must be chosen from the following table, assuming the sizes and alignment given.