

LTL Syntax

Let p range over a given set *Atoms* of atomic propositions.

$$\phi ::= p \mid \neg\phi \mid \phi \wedge \psi \mid \mathbf{X}\phi \mid \mathbf{G}\phi \mid \mathbf{F}\phi \mid \phi \mathbf{U}\psi$$

Models and Paths

A *model* is a tuple $\mathcal{M} = (S, \rightarrow, L)$ where:

- (i) S is a set of states,
- (ii) $\rightarrow \subseteq S \times S$ is a transition relation,
- (iii) $L : S \rightarrow 2^{Atoms}$ is a labelling function.

A *path* of \mathcal{M} is an infinite sequence of states $\pi = s_0s_1s_2s_3\dots$ such that $s_i \rightarrow s_{i+1}$ for all $i \geq 0$. Given such a path, we denote by $\pi(i)$ the i -th element s_i of π , and we denote by π^i the i -th suffix $s_i s_{i+1} s_{i+2} s_{i+3} \dots$ of π .

LTL Semantics

Let $\mathcal{M} = (S, \rightarrow, L)$ be a model, and let π be a path of \mathcal{M} .

$$\begin{aligned} \pi \models^{\mathcal{M}} p &\stackrel{\text{def}}{\iff} p \in L(\pi(0)) \\ \pi \models^{\mathcal{M}} \neg\phi &\stackrel{\text{def}}{\iff} \text{not } \pi \models^{\mathcal{M}} \phi \\ \pi \models^{\mathcal{M}} \phi \wedge \psi &\stackrel{\text{def}}{\iff} \pi \models^{\mathcal{M}} \phi \text{ and } \pi \models^{\mathcal{M}} \psi \\ \pi \models^{\mathcal{M}} \mathbf{X}\phi &\stackrel{\text{def}}{\iff} \pi^1 \models^{\mathcal{M}} \phi \\ \pi \models^{\mathcal{M}} \mathbf{G}\phi &\stackrel{\text{def}}{\iff} \forall i \geq 0. \pi^i \models^{\mathcal{M}} \phi \\ \pi \models^{\mathcal{M}} \mathbf{F}\phi &\stackrel{\text{def}}{\iff} \exists i \geq 0. \pi^i \models^{\mathcal{M}} \phi \\ \pi \models^{\mathcal{M}} \phi \mathbf{U}\psi &\stackrel{\text{def}}{\iff} \exists i \geq 0. (\pi^i \models^{\mathcal{M}} \psi \wedge \forall j < i. \pi^j \models^{\mathcal{M}} \phi) \end{aligned}$$

Model \mathcal{M} at state $s \in S$ satisfies formula ϕ , denoted $\mathcal{M}, s \models \phi$, if all paths π of \mathcal{M} starting at s satisfy ϕ .