

Information Flow Security (2)

DD2460 Software Safety and Security: Part III, lecture 3

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Outline

Information Flow Security deals with *Confidentiality* and *Integrity* related security policies.

- 1 Noninterference Variants
- 2 Enforcement Techniques
- 3 Conclusion / Wrap-up


Noninterference Variants



Termination (In)sensitive Noninterference

Main idea: attacker is (un)able to observe (\mathcal{O}) if execution terminated or not

$$\forall \sigma_1, \sigma_2: \sigma_1 =_L \sigma_2 \Rightarrow \mathcal{O}[[\sigma_1 \vdash P]] = \mathcal{O}[[\sigma_2 \vdash P]]$$

- Sensitive: tag termination into observables
- Insensitive (1): observable prefixes of nonterminating executions
- Insensitive (2): discard non-terminating executions (σ) 

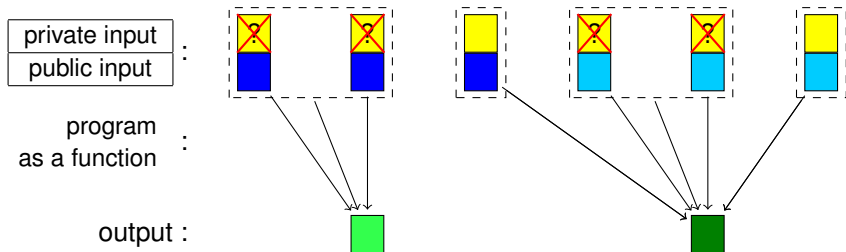


Declassification

Definition 1 (Noninterference modulo declassification ϕ)

A program is safe if and only if any executions, started with the same public inputs *and agreeing on ϕ* , output the same sequence.

$\phi =$ secret is or is not **yellow**





Taint Analysis

Takes into account only (direct) explicit flows

Weaker security guarantees, but more efficient enforcement mechanisms

- not efficient against malicious code, but OK against buggy code

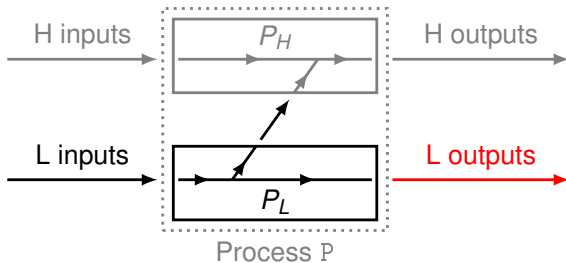
Examples:

- Python's taint library
- Perl taint mode
- ...

Enforcement Techniques



Noninterference Enforcement: Main Idea





Old Security Mechanism: Confined Processes

Lampson's 1973 notion of *confinement*

Confined processes:

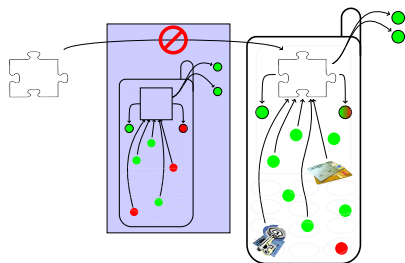
- are memoryless (\Rightarrow side-effect free)
- call only confined processes, but can be called by unconfined processes
- have masked output belonging to a predefined set
 - could extend to label verification

Main concepts underlying *sandboxing*

- one of Java's main security mechanisms



Static Information Flow Analysis



Principles:

- analyze IF before execution
- do nothing during execution

Advantages:

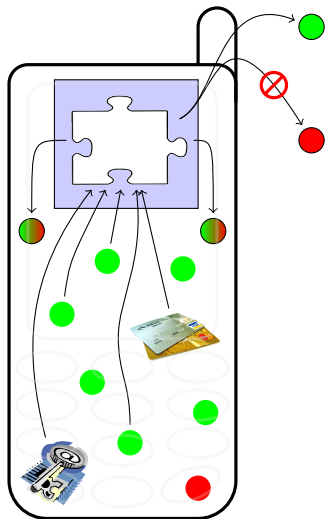
- no runtime overhead
- run iff NI is proved
- old strong soundness culture

Main drawback:

- can be too restrictive



Dynamic Information Flow Analysis



Principles:

- track flows at execution
- prevent data leak just before it occurs

Advantages:

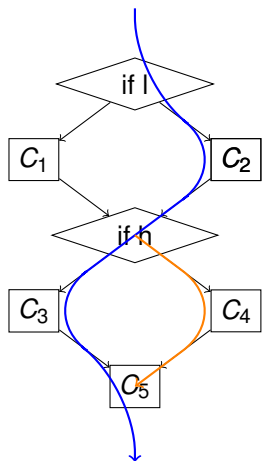
- enforce runtime policies more easily
- allow *safe* executions of unsafe programs
- may be more precise in some cases
 - reduced space (not all executions)
 - access to runtime values

Main drawback:

- hard to spot all flows (implicit flows)



Hybrid Information Flow Analysis



Principles: mix of static and dynamic analyses

- dynamically analyze C_2 and C_3
 - for *direct* and *explicit indirect* flows
- statically analyze C_4
 - for *implicit indirect* flows
- dynamically analyze C_5 with results of C_3 and C_4 mixed

Advantages:

- best of both worlds

Main drawback:

- worst of both worlds
- higher complexity



Is Detection Enough?

What happens with an analysis which is *sound* with regard to information flow detection?

- Static analysis:

Expert: “You should not use this program!”

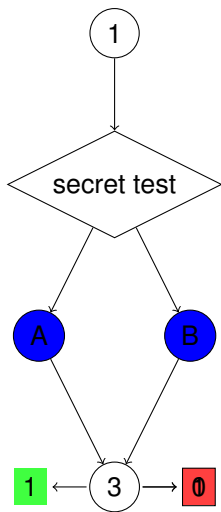
- Dynamic analysis:

ATM: “Oh, by the way, I probably sent your PIN code all over the web.”

A user expects dynamic IF analyses to detect *and correct* information flows.



The Correction Pitfall



Code block A outputs value 1:



Analysis concludes:

- public data: ○ → ■
- secret data: ○ → ■

Sound detection *does not* imply sound (detection + correction)



- dynamic analysis + “stop” correction
- “stop” correction with termination insensitive NI proof

Conclusion / Wrap-up



3 Most Important Points

- \exists many information flow security policy variants
 - termination sensitivity
 - declassification
 - ...
 - taint analyses

- Enforcement
 - Static analyses: (+) soundness (-) usability (often too restrictive)
 - Dynamic analyses: (+) usability (-) soundness
 - Hybrid analyses: (+/-) soundness & usability (-) complexity

- Correction pitfall
 - dynamic and hybrid analyses require correction mechanism
 - sound detection \nRightarrow sound (detection + correction)



IF Workshop

Goal: simulate review of some existing IF security techniques

- you *do not* need to *defend* or *kill* your paper
- you need to:
 - describe the enforcement technique used [and its implementation] (for reproducibility)
 - evaluate the level of security provided
 - describe advantages and limitations of the technique
 - compare with other known techniques:
 - workshop: type system + taint analysis
 - report: type system + taint analysis + workshop techniques

After the workshop and report, I/you should be able to pick up the best adapted tool/technique for a particular IF problem.



Grading

Workshop presentation is not graded per se (report is)

[due 12/3]

- E:
 - give a decent presentation (or at least additions/corrections session)
 - be able to give an accurate description/summary of the paper at the course level
- C: (subsumes E)
 - detail specific advantages and limitations of the paper's technique
- A: (subsumes A)
 - compare with the relevant techniques presented in class and in the other papers

Level of learning of course material also reflected in the final grade

- if/where possible, report should contain proof of knowledge of channels, flows, labels, noninterference, enforcement, . . .



Information Flow Wrap-up

		Enforcement		
		Type System	Taint	Others
basic	Concepts Definitions lectures 1 & 3: IF policies, channels, flows, labels, correction, ... exercises 1 & 2: IF policies, timing channels, flows, ...	lecture 2: type systems, noninterference, ... exercises 2: type systems, Jif, ...	exercises 1: taint, ...	lecture 3: static, dynamic, hybrid, ... workshop
	deeper			

Annotations: A red oval highlights the 'basic' row. An orange oval highlights the 'Concepts Definitions' column and the 'Type System' column. A green rectangle highlights the entire table content.



Course Wrap-up

Software safety and security:

- prevent bad behaviors causing system (base) and data (load) damage
- due to specification and/or implementation errors and/or weaknesses

Formal methods:

- precise correctness guarantees
- often complex and expensive
- for critical systems and/or data

3 different techniques for software safety and security

- Temporal logic and model checking
- Hoare logic and VCG/symbolic execution
- Information flow and type system



Announcements and Questions?

Soon online:

- lab 2 booking
- course evaluation

Questions?