Passwords in Peer-to-Peer

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Need for Logins P2P Solution

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Modern End-User Services









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Logging in to P2P?

What was the last P2P application you logged in to?



Background Basic Solution More Features

Evaluation and Conclusion

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The Easy Case

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Need for Logins P2P Solution

Sometimes we Want Authentication





Need for Logins P2P Solution

Standard P2P Solution

- Generate a cryptographic key
- Whoever knows key is authenticated as user
- On a single machine it Just WorksTM



Need for Logins P2P Solution

Standard P2P Solution

- Generate a cryptographic key
- Whoever knows key is authenticated as user
- On a single machine it Just WorksTM
- but, needs to be copied to smartphone.



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Standard P2P Solution Depicted





Need for Logins P2P Solution

Problem Crops up in P2P Backup

- Default solution fails spectacularly for backup
- Some punt problem to user: back up the key separately
- Some derive master key from a password



Target Scenario

Application



Lo	gin Proto	col
DHT	Dist. Storage	Peer Sampling



Target Scenario

Application





Target Scenario





The Basic Solution

- ► Basic solution: derive key from user-entered password
- Cryptographic key-derivation function (PBKDF-2, bcrypt)



Basic Solution Depicted





Adding some flexibility

- Tying key we use to password can be limiting
- Add a key store where keys we need are stored
- Key store is encrypted with key derived from password



Locating the key store

- How does one find the key store?
- Lookup service (DHT) mapping username to storage location
- Needs to be write-once to prevent DoS



Flexible Basic Solution Depicted





Why not just use a Cloud?

- Censorship resistance
- Administrative cost
- Requires trust in cloud provider



Why not just use a Cloud?

- Censorship resistance
- Administrative cost
- Requires trust in cloud provider
- Our protocol can be deployed on cloud storage



Remember Me Password Change Forgotten Passwords

Remember Me





Remember Me Password Change Forgotten Passwords

Requirements

- Password must not be recoverable from device
- Must be able to revoke stored credentials on other devices



Remember Me Password Change Forgotten Passwords

Design

 All problems in computer science can be solved by another level of indirection



Remember Me Password Change Forgotten Passwords

Design

- All problems in computer science can be solved by another level of indirection
- For each device, a device login information file in distributed storage
- ► File contains encrypted copy of key store key
- Device stores key to decrypt its device login file



Remember Me Password Change Forgotten Passwords

Simplified Remember Me in Pictures





Remember Me Password Change Forgotten Passwords

Revokation

- To support revokation, devices' keys are also escrowed
- Key store key is replaced
- All non-revoked devices' login information files are updated
- Application keys from key store must also be updated!



Remember Me Password Change Forgotten Passwords

Remember Me in Pictures





Remember Me Password Change Forgotten Passwords

Remember Me in Pictures





Remember Me Password Change Forgotten Passwords

Remember Me in Pictures





Remember Me Password Change Forgotten Passwords

Password Change

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New e	nail addre	SS		
New p	assword			
Please (change:	enter your o	old passwo	ord to confirn	n your
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Remember Me Password Change Forgotten Passwords

Requirements

User must know old password to change it



Remember Me Password Change Forgotten Passwords

Authentication for distributed storage

How to authenticate for the distributed storage (Catch-22)?



Remember Me Password Change Forgotten Passwords

Authentication for distributed storage

- How to authenticate for the distributed storage (Catch-22)?
- With cryptographic keys, of course! (It's a normal P2P protocol)
- All files are world-readable, only writing requires key
- Write-permission for most files from key store
- Write-permission for login info only if one can decrypt login info



Remember Me Password Change Forgotten Passwords

Design

- Using password, retrieve key needed to update login info
- Generate new key for key store
- Store new key store
- Update login info to point to new key store
- Update/remove device login information files



Remember Me Password Change Forgotten Passwords

Forgotten Passwords





Remember Me Password Change Forgotten Passwords

Requirements

- Should not reveal secret information to others
- Based on e-mail and/or security questions



Remember Me Password Change Forgotten Passwords

Design

- Two separate schemes for e-mail and questions
- Building on secret sharing
- For questions, shares are encrypted with key derived from answers
- ► For e-mail, shares are spread among random peers
- Using cryptographic commitments to only reveal e-mail address when actually used



Evaluation Setting

- Simulation using timing from Bittorrent DHT
- DHT timing by Jiménez et al (P2P'11)
- Excludes crypto computation (negligible aside from KDF)



Evaluation Summary

Evaluation Results





Evaluation Summary

Future Work

- Security proofs
- Prevent offline guessing attacks
- Implementation





- Fully distributed protocol supporting main features of password-based login
- Including "Remember me", password change, password reset
- Room for future improvements

