

Lecture 1

Course Description and Introduction to Ciphers

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Introduction and Administration

What is cryptography?

Cryptography is concerned with the conceptualization, definition, and construction of computing systems that address security concerns.

– Oded Goldreich, Foundations of Cryptography, 1997

Applications of Cryptography

Historically.

- ▶ Military and diplomatic secret communication.
- ▶ Communication between banks, e.g., credit card transactions.

Modern Time.

- ▶ Protecting satellite TV from leaching.
- ▶ Secrecy and authenticity on the Internet, mobile phones, etc.
- ▶ Credit cards.

Applications of Cryptography

Today.

- ▶ Distributed file systems, authenticity of blocks in bit torrents, anonymous remailers, Tor-network, etc.
- ▶ RFID tags, Internet banking, Försäkringskassan, Skatteverket, “e-legitimation”.

Future.

- ▶ Secure distributed computing (multiparty computation): election schemes, auctions, secure cloud computing, etc.
- ▶ Variations of signatures, cryptosystem, and other primitives with special properties, e.g., group signatures, identity based encryption, etc.

Goal

The goal of the course is to

- ▶ give an overview of modern cryptography

in order that students should

- ▶ know how to evaluate and, to some extent, create cryptographic constructions, and
- ▶ to be able to read and to extract useful information from research papers in cryptography.

Prerequisites

- ▶ *DD1352 Algorithms, data structures and complexity, or DD2354 Algorithms and complexity.*
- ▶ Knowledge of mathematics and theory of algorithms corresponding to the required courses of the D or F-programmes at KTH.

Tentative Plan of Content (1/2)

- ▶ Administration, introduction, classical cryptography.
- ▶ Symmetric ciphers, substitution-permutation networks, linear cryptanalysis, differential cryptanalysis.
- ▶ AES, Feistel networks, Luby-Rackoff, DES, modes of operations, DES-variants.
- ▶ Entropy and perfect secrecy.
- ▶ Repetition of elementary number theory,
- ▶ Public-key cryptography, RSA, primality testing, textbook RSA, semantic security.

Tentative Plan of Content (2/2)

- ▶ RSA in ROM, Rabin, discrete logarithms, Diffie-Hellman, El Gamal.
- ▶ Security notions of hashfunctions, random oracles, iterated constructions, SHA, universal hashfunctions.
- ▶ Message authentication codes, identification schemes, signature schemes, PKI.
- ▶ Elliptic curve cryptography.
- ▶ Pseudorandom generators.
- ▶ Guest lecture.
- ▶ Make-up time and/or special topic.

Course Requirements

Presentations. **a)** Choose a research topic, and **b)** summarize the paper in a 12-min oral presentation.

Gives P -points ($P = 0$ or $30 \leq P \leq 80$), which is the sum of:

- ▶ (20P) Choice of content.
- ▶ (20P) Understanding of the content
- ▶ (20P) Quality of slides (or whiteboard)
- ▶ (20P) Presentation skills.

Up to 6 talks in 2 hour-sessions. Listen to the talks in your session.

Detailed rules and advice:

<http://www.csc.kth.se/DD2448/krypto14/handouts/talk.pdf>

Course Requirements

Homework 1-4. Each homework is a set of problems giving I -points and T -points ($I \geq 10$ and $I + T \geq 50$).

- ▶ Solved in groups of up to three students, which may differ for each homework.
- ▶ Only informal discussions are allowed.
- ▶ Each student writes and submits his own solution.

Detailed rules and advice:

http://www.csc.kth.se/DD2448/krypto14/handouts/solution_rules.pdf

Course Requirements

Oral Exam. Purpose is to give a fair grade.

Discussion starting from submitted solutions and presentation.

Gives (possibly negative) I -points and T -points and a single O -point if passed.

Grading

To earn a given grade the requirements of all lower grades must be satisfied as well, with $A = I + T + P + O$.

Grade	Requirements
E	$I \geq 30$, $T \geq 40$, $P \geq 30$, and $O \geq 1$.
D	$A \geq 120$.
C	$A \geq 140$ and $P \geq 50$.
B	$A \geq 170$.
A	$A \geq 210$ and $P \geq 60$.

Kattis

Kattis is a judging server for programming competitions and for grading programming assignments. We use it for all exercises where code is submitted as a solution.

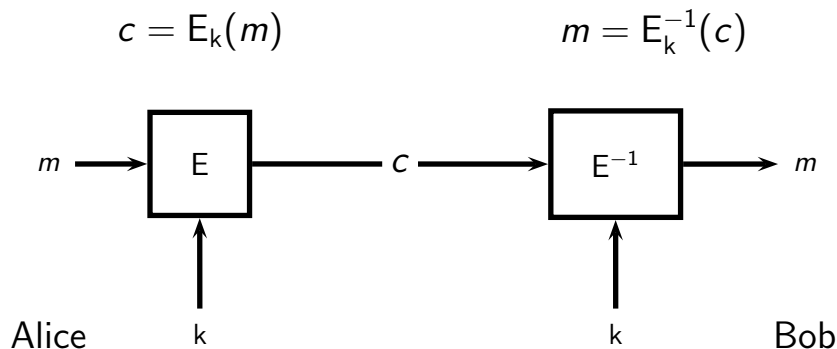
We assume that your Kattis id is the same as your KTH user name. If this is not the case, then email us your Kattis user name and use the subject Krypto14 Kattis.

Latex

- ▶ Latex is the standard typesetting tool for mathematics.
- ▶ It is the fastest way to produce mathematical writing. You must use it to typeset your solutions.
- ▶ The best way to learn it is to read:
<http://tobi.oetiker.ch/lshort/lshort.pdf>

Introduction to Ciphers

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such that $E_k^{-1}(E_k(m)) = m$ for every message $m \in \mathcal{P}$ and $k \in \mathcal{K}$. The set $\mathcal{C} = \{E_k(m) \mid m \in \mathcal{P} \wedge k \in \mathcal{K}\}$ called the **set of ciphertexts**.

Attacks

Throughout the course we consider various attacks on cryptosystems. With small changes, these attacks make sense both for symmetric and asymmetric cryptosystems.

- ▶ Ciphertext-only attack.
- ▶ Known-plaintext attack
- ▶ Chosen-plaintext attack
- ▶ Chosen-ciphertext attack

Cesar Cipher (Shift Cipher)

Consider English, with alphabet A-Z_, where _ denotes space, thought of as integers 0-26, i.e., \mathbb{Z}_{27}

- ▶ **Key.** Random letter $k \in \mathbb{Z}_{27}$.
- ▶ **Encrypt.** Plaintext $m = (m_1, \dots, m_n) \in \mathbb{Z}_{27}^n$ gives ciphertext $c = (c_1, \dots, c_n)$, where $c_i = m_i + k \pmod{27}$.
- ▶ **Decrypt.** Ciphertext $c = (c_1, \dots, c_n) \in \mathbb{Z}_{27}^n$ gives plaintext $m = (m_1, \dots, m_n)$, where $m_i = c_i - k \pmod{27}$.

Cesar Cipher Example

Encoding.

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	_
00	01	02	03	04	05	06	07	08	09	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26

Key: $G = 6$

Plaintext. B R I B E _ L U L A _ T O _ B U Y _ J A S

Plaintext. 01 17 08 01 04 26 11 20 11 00 26 19 14 26 01 20 24 26 09 00 18

Ciphertext. 07 23 14 07 10 05 17 26 17 06 05 25 20 05 07 26 03 05 15 06 24

Ciphertext. H X O H K F R _ R G F Z U F H _ D F P G Y