

## Lecture 3: Linear Algebra: Minimization and equilibrium, S. Ch 1 & 2

The convergence of a Markov chain density function to a steady state is easy to show for diagonalizable transition matrices  $\mathbf{W}$ . What about non-diagonalizable  $\mathbf{W}$ ? This question is of more general interest for dynamical systems and iterative solution of equations. So, consider the powers  $\mathbf{A}^n$ ,  $n = 1, 2, \dots$  of a  $m \times m$  real matrix  $\mathbf{A}$ .

Definition: The spectral radius is the maximal modulus of any eigenvalue,

$$\rho(\mathbf{A}) = \max |\lambda_i|$$

Theorem. If  $\rho(\mathbf{A}) < 1$ ,  $\lim_{n \rightarrow \infty} \mathbf{A}^n = \mathbf{0}$

We will outline the proof, leaving some details out. First, the *Schur* theorem guarantees that any square matrix can be triangularized by a unitary matrix  $\mathbf{Q}$ :

$$\mathbf{A} = \mathbf{Q}\mathbf{U}\mathbf{Q}^H, \mathbf{Q}\mathbf{Q}^H = \mathbf{I}$$

This is a similarity transformation:  $\mathbf{U}$  and  $\mathbf{A}$  have the same eigenvalues.  $\mathbf{U}$  is upper triangular, and  $\mathbf{Q}$  can be chosen to put the eigenvalues of  $\mathbf{A}$  in any order on the diagonal of  $\mathbf{U}$ . The proof of this relies on the fact that any matrix has an eigenvalue and an eigenvector, but does not tell how to calculate it.

So, let there be  $q$  different eigenvalues, and arrange them in blocks down the diagonal,

$$\mathbf{U} = \begin{pmatrix} \mathbf{U}_1 & x & x & x \\ 0 & \mathbf{U}_2 & x & x \\ \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \mathbf{U}_q \end{pmatrix}, \mathbf{U}_k = \lambda_k (\mathbf{I} + \mathbf{N}_k), k = 1, \dots, q,$$

$$\mathbf{N}_k = \begin{pmatrix} 0 & x & x & \dots & x \\ 0 & 0 & x & \dots & x \\ \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & \dots & 0 & x \\ 0 & 0 & \dots & 0 & 0 \end{pmatrix}, n_k \times n_k$$

$\mathbf{N}_k$  is upper triangular with zeros on the diagonal, and so *nilpotent*,  $\mathbf{N}_k^{n_k} = \mathbf{0}$ .

One can also find a similarity transformation  $\mathbf{S}$  (but not unitary) such that

$$\mathbf{S}\mathbf{U}\mathbf{S}^{-1} = \begin{pmatrix} \mathbf{U}_1 & 0 & \dots & 0 \\ 0 & \mathbf{U}_2 & \dots & 0 \\ \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \mathbf{U}_q \end{pmatrix} = \mathbf{F}, \mathbf{F}^n = \begin{pmatrix} \mathbf{U}_1^n & 0 & \dots & 0 \\ 0 & \mathbf{U}_2^n & \dots & 0 \\ \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \mathbf{U}_q^n \end{pmatrix}$$

so we can now focus attention on the powers of the diagonal blocks  $\mathbf{U}_k$

The binomial expansion says, for an  $m \times m$  matrix  $\mathbf{U}$ ,

$$\mathbf{U}^n = \lambda^n (\mathbf{I} + \mathbf{N})^n = \lambda^n \sum_{k=0}^n \binom{n}{k} \mathbf{N}^k = \lambda^n \sum_{k=0}^{m-1} \binom{n}{k} \mathbf{N}^k$$

because the unit matrix commutes with any matrix, and all powers  $> m-1$  of  $\mathbf{N}$  vanish.

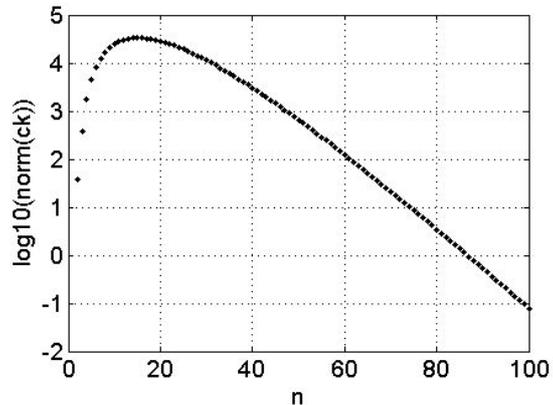
The final step uses a norm estimate,

$$\|U^n\| \leq |\lambda|^n \sum_{k=0}^{m-1} \binom{n}{k} \|N\|^k \leq |\lambda|^n \max(m, \|N\|^{m-1}) n^{m-1} = C |\lambda|^n n^{m-1}$$

which tends to 0 as  $n$  grows when  $|\lambda| < 1$ . This finishes the proof.

However, the matrix grows polynomially initially. Here is an example:

$$U = \begin{pmatrix} a & m & m & m \\ 0 & a & m & m \\ 0 & 0 & a & m \\ 0 & 0 & 0 & a \end{pmatrix}, a = 0.8, m = 10$$



The plot shows  $\|U^n \mathbf{c}\|_2$  vs.  $n$ ,  $\mathbf{c} = (1, 1, 1, 1)^T$ .

**Least squares approximation: Normal equations, QR, and SVD.**

Ex. Given data points  $(x_i, f_i)$ ,  $i = 1, 2, \dots, m$ , find a polynomial  $p(x) = a_0 + a_1x + a_2x^2$  which approximates the data,  $p(x_i) = f_i$ . This is a linear system  $\mathbf{V}\mathbf{a} = \mathbf{f}$ ,

$$\begin{pmatrix} 1 & x_1 & x_1^2 \\ 1 & x_2 & x_2^2 \\ 1 & x_3 & x_3^2 \end{pmatrix} \begin{pmatrix} a_0 \\ a_1 \\ a_2 \end{pmatrix} = \begin{pmatrix} f_1 \\ f_2 \\ f_3 \end{pmatrix}$$

for  $m = 3$ : When the  $x_i$  are distinct, there is a unique interpolation polynomial for any data  $(x_i, f_i)$ . It may not be obvious that the columns of  $\mathbf{V}$  are linearly independent, but we may compute the polynomial by another ansatz:

$$\begin{aligned} p(x) &= c_0 + c_1(x - x_1) + c_2(x - x_1)(x - x_2) : \\ f_1 &= p(x_1) = c_0 + 0 + 0 \\ f_2 &= p(x_2) = c_0 + c_1(x_2 - x_1) + 0 \\ f_3 &= p(x_3) = c_0 + c_1(x_3 - x_1) + c_2(x_3 - x_1)(x_3 - x_2) \end{aligned}$$

$$\begin{pmatrix} 1 & 0 & 0 \\ 1 & x_2 - x_1 & 0 \\ 1 & x_3 - x_1 & (x_3 - x_1)(x_3 - x_2) \end{pmatrix} \begin{pmatrix} c_0 \\ c_1 \\ c_2 \end{pmatrix} = \begin{pmatrix} f_1 \\ f_2 \\ f_3 \end{pmatrix}$$

The coefficient matrix is lower triangular, and the system can be solved for  $\mathbf{c}$  as long as the diagonal elements are non-zero. But the  $c$ -form and the  $a$ -form both generate all quadratic polynomials, so this shows that the system for the  $a$ -form is always non-singular. Indeed, the *Vandermonde* determinant may be calculated:

$$\det \begin{pmatrix} 1 & x_1 & x_1^2 \\ 1 & x_2 & x_2^2 \\ 1 & x_3 & x_3^2 \end{pmatrix} = (x_3 - x_1)(x_3 - x_2)(x_2 - x_1)$$

Next, we consider  $m > 3$ . We choose to find coefficients to minimize the sum of squares of discrepancies,

$$\min_a \sum_{j=1}^m (f_j - p(x_j))^2 = \min_a (\mathbf{r}, \mathbf{r}), \mathbf{r} = (r_1, r_2, \dots, r_m)^T, r_j = f_j - p(x_j)$$

The development exploits the scalar product (inner-product) ( $\cdot, \cdot$ ),

**Ex.**

For the vector space  $\mathbf{R}^n$  of real  $n$ -vectors, the standard inner product is

$$(\mathbf{x}, \mathbf{y}) = \sum x_i y_i = \mathbf{x}^T \mathbf{y}$$

and we can define the Euclidean vector norm  $\|\mathbf{x}\|_2^2 = (\mathbf{x}, \mathbf{x})$  with the Cauchy-Schwarz inequality  $(\mathbf{x}, \mathbf{y}) \leq \|\mathbf{x}\|_2 \|\mathbf{y}\|_2$ . So we can define angles between vectors,

$$\frac{(\mathbf{x}, \mathbf{y})}{\|\mathbf{x}\|_2 \|\mathbf{y}\|_2} = \cos \theta, \text{ etc.}$$

$\mathbf{x}$  and  $\mathbf{y}$  are orthogonal if  $(\mathbf{x}, \mathbf{y}) = 0$ .

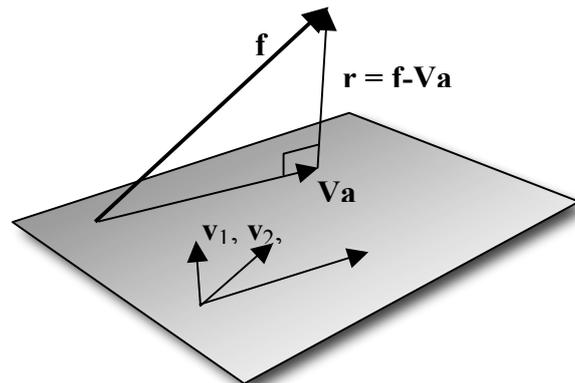
The optimum coefficients  $\mathbf{a}$  - giving residuals  $\mathbf{r}^*$  - is characterized by:

Let  $\mathbf{r} = \mathbf{f} - \sum a_j \mathbf{v}_j$  (written  $\mathbf{r} = \mathbf{f} - \mathbf{V}\mathbf{a}$  above)

Then  $(\mathbf{r}^*, \mathbf{r}^*) \leq (\mathbf{r}, \mathbf{r})$  if and only if  $(\mathbf{r}^*, \mathbf{v}) = 0$  for all  $\mathbf{v}$  in the column space of  $\mathbf{V}$ , i.e., *the optimal residual vector is normal to all vectors in  $\mathbf{V}$  - the normal equations.*

Here is the picture:

The point  $\mathbf{V}\mathbf{a}^*$  in the subspace spanned by the  $\mathbf{v}_i$  has minimal distance to  $\mathbf{f}$ . Perturbing the  $\mathbf{a}$  to  $\mathbf{a}^* + s\mathbf{c}$  changes  $\mathbf{V}\mathbf{a}$  by  $s\mathbf{v} = s\mathbf{V}\mathbf{c}$  ( $s$  is a scalar multiple) and  $\mathbf{r}$  by  $-s\mathbf{v}$ .



So for all  $s$ ,

$$\begin{aligned} (\mathbf{r}^*, \mathbf{r}^*) &\leq (\mathbf{r}^* - s\mathbf{v}, \mathbf{r}^* - s\mathbf{v}) = (\mathbf{r}^*, \mathbf{r}^*) + s^2 (\mathbf{v}, \mathbf{v}) - 2s(\mathbf{r}^*, \mathbf{v}) = \\ &= (\mathbf{r}^*, \mathbf{r}^*) + (\mathbf{v}, \mathbf{v}) \left( s - \frac{(\mathbf{r}^*, \mathbf{v})}{(\mathbf{v}, \mathbf{v})} \right)^2 - \frac{(\mathbf{r}^*, \mathbf{v})^2}{(\mathbf{v}, \mathbf{v})}. \end{aligned}$$

Choosing  $s = (\mathbf{r}^*, \mathbf{v}) / (\mathbf{v}, \mathbf{v})$  we see that only  $(\mathbf{r}^*, \mathbf{v}) = 0$  for all  $\mathbf{v}$  in the subspace can satisfy the inequality.

We obtain the normal equations by taking  $\mathbf{v} = \mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_m$ , the columns of  $\mathbf{V}$ :

$$\mathbf{V}^T \mathbf{V}\mathbf{a} = \mathbf{V}^T \mathbf{f}$$

The system can be solved by **LU** or **LDL<sup>T</sup>** factorization. We will look at another idea: To obtain an orthogonal basis for **V** by e.g. the Gram-Schmidt orthogonalization.

Suppose for the moment that the  $v_i$  are orthogonal, let us call them  $q_i$ . Then

$$\mathbf{Q}^T \mathbf{Q} \mathbf{a} = \mathbf{Q}^T \mathbf{f}, \text{ and } \mathbf{Q}^T \mathbf{Q} = \text{diag}(\mathbf{q}_k^T \mathbf{q}_k), \text{ so}$$

$$a_k = \mathbf{q}_k^T \mathbf{f} / \mathbf{q}_k^T \mathbf{q}_k, k = 1, 2, \dots, m$$

Here is a variant of the Gram-Schmidt algorithm:

$$q_1 := v_1 / \|v_1\|$$

for  $k = 2, 3, \dots$

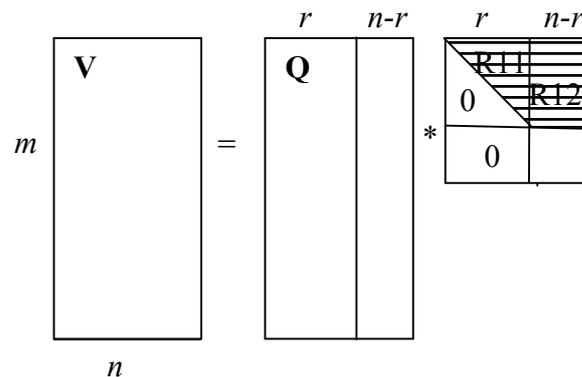
$$q_k := v_k - \sum_{j=1}^{k-1} r_{jk} q_j; r_{jk} = (v_k, q_j)$$

$$q_k := q_k / \|q_k\|$$

end

The values of  $r_{jk}$  makes  $q_k$  orthogonal to all earlier  $q_j$ , and they are normalized to unit length. If the  $v_j$  are linearly dependent,  $q_k$  may become zero before one has used all the  $v_j$ . The algorithm is therefore combined with column reordering to choose the largest remaining  $v$  at every step.

Then, the process finishes with an upper triangular **R**-matrix whose last  $n-r$  rows are zeros, where  $r$  is the column rank of **V**. The corresponding columns of **Q** can be chosen arbitrarily, orthogonal to the  $r$  first.



From  $\mathbf{V} \mathbf{P} = \mathbf{Q} \mathbf{R}$  follows  $\mathbf{Q} \mathbf{R} \mathbf{P}^T \mathbf{a} = \mathbf{f}$   
 so with  $\mathbf{R} \mathbf{P}^T \mathbf{a} = \mathbf{y}$  we have

$$y_k = 0, k = r+1, r+2, \dots, n$$

and the solution to the normal equations

$$y_i = (\mathbf{q}_i, \mathbf{f}), i = 1, 2, \dots, r$$

which gives the minimal distance. If  $r < n$ , we get

$$\mathbf{a}_1 = \mathbf{R}11^{-1} (\mathbf{y}_1 - \mathbf{R}12 \mathbf{a}_2)$$

where  $\mathbf{y}_1 = (y_1, y_2, \dots, y_r)^T$ , etc.

A unique **a**-solution can be defined as the one with minimal number of non-zeros, i.e.,  $\mathbf{a}_2 = 0$ . This is what Matlab's backslash gives:

$$\mathbf{a} = \mathbf{V} \backslash \mathbf{f};$$

Choosing the solution of minimal l2-norm defines the pseudo-inverse,  $\mathbf{V}^+$ ,

$$\mathbf{a} = \text{pinv}(\mathbf{V}) * \mathbf{f};$$

This is computed by the singular value decomposition, developed into a practical tool by the Gene Golub (-2007) and Cleve Moler.

Any real  $m \times n$  matrix  $\mathbf{A}$  admits the factorization

$$\mathbf{A} = \mathbf{U}\mathbf{S}\mathbf{V}^T,$$

where  $\mathbf{U}$  is  $m \times m$ ,  $\mathbf{V}$  is  $n \times n$ , both orthogonal. The first  $r$  columns of  $\mathbf{U}$  is an orthogonal basis for the column space of  $\mathbf{A}$ , and the  $r$  first columns of  $\mathbf{V}$  are an orthogonal basis for the row space.  $\mathbf{S}$  is  $m \times n$ , non-zeros only on the diagonal  $s_{ii} = \sigma_i$ , sorted  $\sigma_1 > \sigma_2 > \dots > \sigma_r$ , the singular values of  $\mathbf{A}$ .

The  $m-r$  last columns of  $\mathbf{U}$  can be chosen at will, if orthogonal to the  $r$  first columns, d:o for  $\mathbf{V}$ .

The pseudo-inverse  $\mathbf{S}^+$  is obtained by inverting the non-zeros of  $\mathbf{S}$ , so

$$\mathbf{A}^+ = \mathbf{V}\mathbf{S}^+\mathbf{U}^T$$

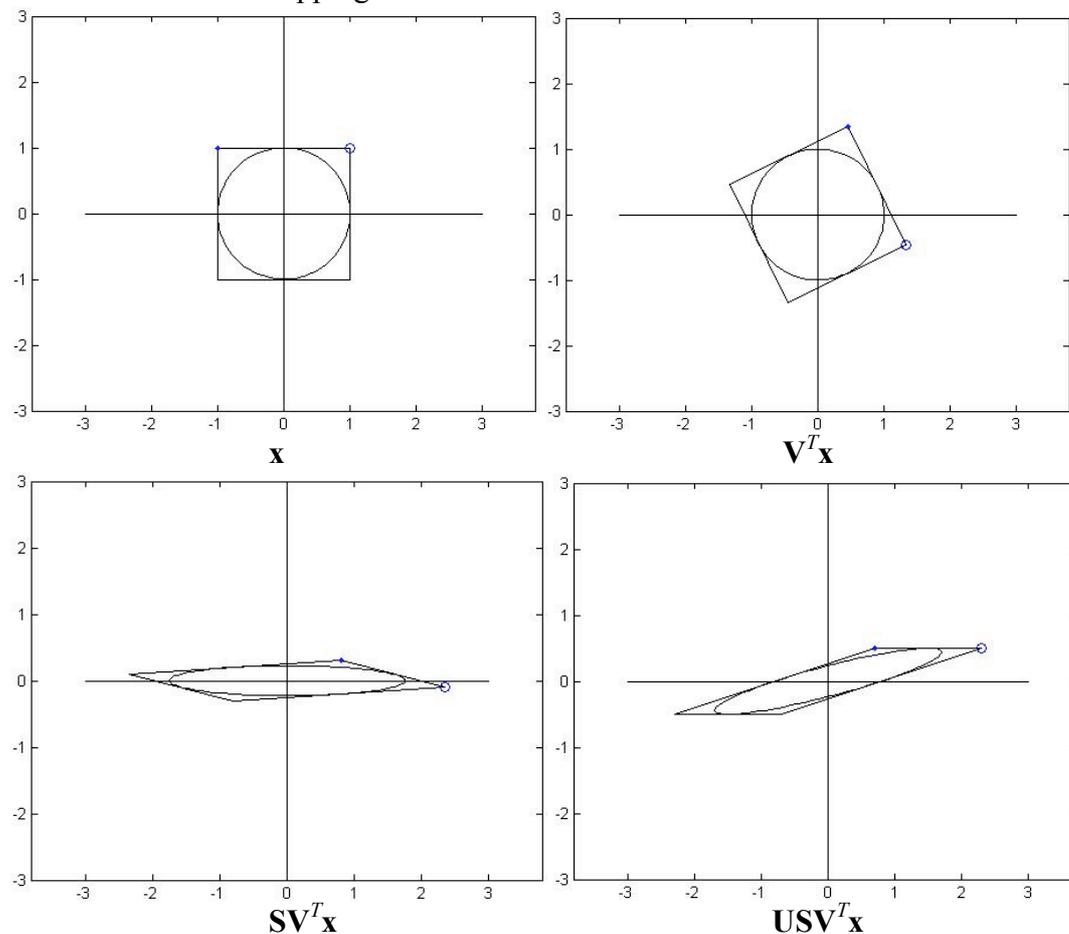
The SVD can in principle be computed from eigenvalues and -vectors of  $\mathbf{A}^T\mathbf{A}$  but the Golub-Reinsch algorithm uses a bi-diagonalization procedure which avoids the formation of the matrix product.

**Ex.** What is the SVD of an  $m \times n$  rank-1 matrix  $\mathbf{A} = \mathbf{u}\mathbf{v}^T$

The only non-zero singular value is  $\|\mathbf{u}\| \|\mathbf{v}\|$ , the first column of  $\mathbf{U}$  is  $\mathbf{u}/\|\mathbf{u}\|$ , d:o  $\mathbf{V}$ .

The rest of  $\mathbf{U}$  (and  $\mathbf{V}$ ) is “arbitrary” and can be computed by orthogonalizing a set of linearly independent vectors (such as the set of unit vectors) against  $\mathbf{u}$ , etc.

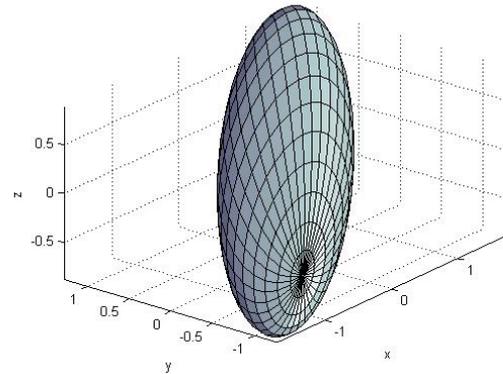
The Singular Value Decomposition describes a linear mapping as a rotation (possibly with a reflection)), followed by a stretching of the coordinate axes, and another rotation. Here is a mapping  $R^2 \rightarrow R^2$



$$\mathbf{A} = \begin{pmatrix} 0.8 & 1.5 \\ 0 & 0.5 \end{pmatrix} = \mathbf{U}\mathbf{S}\mathbf{V}^T, \mathbf{S} = \begin{pmatrix} 1.7573 & 0 \\ 0 & 0.2276 \end{pmatrix},$$
$$\mathbf{V} = \begin{pmatrix} 0.4401 & -0.8979 \\ 0.8979 & 0.4401 \end{pmatrix}, \phi = 64^\circ, \mathbf{U} = \begin{pmatrix} 0.9668 & -0.2555 \\ 0.2555 & 0.9668 \end{pmatrix}, \phi = 15^\circ$$

The singular values are the half-axes of the ellipsoidal image of the unit ball:

$$\mathbf{A} = \begin{pmatrix} 1 & 1 & 1 \\ 0 & 0.8 & 1 \\ 0.2 & 0.7 & 0.5 \end{pmatrix},$$
$$\sigma = 2.2311, 0.6397, 0.1822$$



**Mechanical models: Balls on springs.**

We consider *Hookean* springs, for which the restoring force is proportional to the extension/compression of the spring:

$$F = -K(l - l_0)$$

where  $l$  is the extended length and  $l_0$  is called the natural (force-free) length.

A torsion spring produces a torque proportional to the rotation angle,

$$M = -K \phi$$

The work done *against the spring force* when the spring is extended from length  $l_0$  to  $l_0 + e$  is

$$W = \int_0^e Kl \cdot dl = 1/2 Kl^2$$

so, there is energy stored in the process,  $W = 1/2 Ke^2$ . It follows, that

$$F = - \frac{dW}{dl}$$

Ex. Homogeneous gravity field directed in the negative  $z$ -direction, on a mass point  $m$

$$\mathbf{F} = -gm\mathbf{e}_z$$

The work an external force has to do to move from A to B is

$$W = \int_A^B gm\mathbf{e}_z \cdot d\mathbf{r} = gm \int_A^B dz = mg(z(B) - z(A))$$

The work is independent of the path between A and B. It follows, that the force is the negative gradient of the *potential energy* function,

$$\mathbf{F} = - \frac{\partial W}{\partial z} \mathbf{e}_z$$