From Small Space to Small Width in Resolution

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FLoC Workshop on Proof Complexity Vienna, Austria 13 July 2014

Joint work with Yuval Filmus, Massimo Lauria, Jakob Nordström, and Marc Vinyals

Resolution

Resolution

• Input: CNF formula F

$$(x \vee \overline{y} \vee z) \wedge (\overline{y} \vee \overline{z}) \wedge (x \vee y) \wedge (\overline{x} \vee \overline{z}) \wedge (\overline{x} \vee z)$$

Resolution rule:

$$\frac{C \vee x \qquad D \vee \overline{x}}{C \vee D}$$

ullet Goal: Proof of unsatisfiability (refutation) = Derive empty clause ot

Refer to clauses of formula as axioms

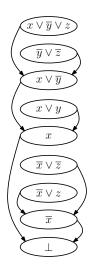
Can represent refutation as

- annotated list or
- DAG

1.	$x \vee \overline{y} \vee z$	Axiom
2.	$\overline{y} \vee \overline{z}$	Axiom
3.	$x \vee \overline{y}$	Res(1,2)
4.	$x \vee y$	Axiom
5.	x	Res(3,4)
6.	$\overline{x} \vee \overline{z}$	Axiom
7.	$\overline{x} \lor z$	Axiom
8.	\overline{x}	Res(6,7)
9.	\perp	Res(5,8)

Can represent refutation as

- annotated list or
- DAG



Can represent refutation as

- annotated list or
- DAG

Size: number of steps in refutation

Space: memory usage (at step t:

clauses before t used after t)

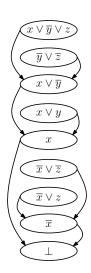
Width: size of the largest clause

Example:

Size

Space

Width



Can represent refutation as

- annotated list or
- DAG

Size: number of steps in refutation

Space: memory usage (at step t:

clauses before t used after t)

Width: size of the largest clause

Example:

Size 9

Space Width 1.

 $x \vee \overline{y} \vee z$

Axiom

 $\overline{y} \vee \overline{z}$

Axiom

3.

 $x \vee \overline{y}$

Res(1,2)

4.

 $x \vee y$

Axiom

5.

x

Res(3,4)

6.

Axiom

7.

 $\overline{x} \vee z$

 $\overline{x} \vee \overline{z}$

Axiom

8.

 \overline{x}

Res(6,7)

9.

Res(5, 8)

Can represent refutation as

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Size: number of steps in refutation

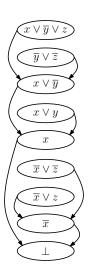
Space: memory usage (at step t:

clauses before t used after t)

Width: size of the largest clause

Example:

Size Space Width



Can represent refutation as

- annotated list or
- DAG

Size: number of steps in refutation

Space: memory usage (at step *t*:

clauses before t used after t)

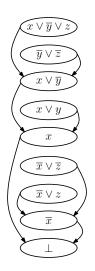
Width: size of the largest clause

Example:

Size S

Space (

Width



Space at current step

Can represent refutation as

- annotated list or
- DAG

```
Size: number of steps in refutation
```

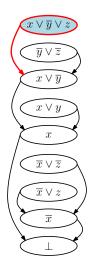
```
Space: memory usage (at step t:
```

clauses before t used after t)

Width: size of the largest clause

Example:

Size 9 Space 1 Width



Space at current step 1

Can represent refutation as

- annotated list or
- DAG

```
Size: number of steps in refutation
```

```
Space: memory usage (at step t:
```

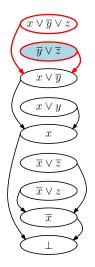
clauses before t used after t)

Width: size of the largest clause

Example:

Size 9 Space 2

Width



Space at current step

Can represent refutation as

- annotated list or
- DAG

Size: number of steps in refutation

Space: memory usage (at step t:

clauses before t used after t)

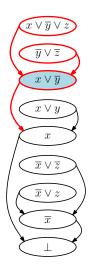
Width: size of the largest clause

Example:

Size 9

Space 3

Width



Space at current step

Can represent refutation as

- annotated list or
- DAG

Size: number of steps in refutation

Space: memory usage (at step t: # clauses before t used after t)

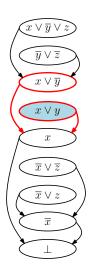
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Size 9

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Space at current step

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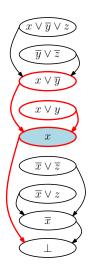
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Space 3

Width



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Space: memory usage (at step t:

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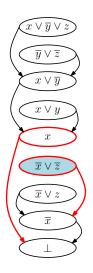
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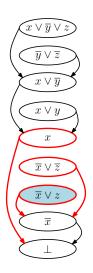
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Space at current step

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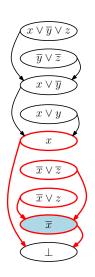
Width: size of the largest clause

Example:

Size 9

Space

Width



Space at current step

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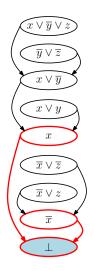
clauses before t used after t)

Width: size of the largest clause

Example:

Size 9 Space 4

Width



Space at current step

Can represent refutation as

- annotated list or
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Size: number of steps in refutation

Space: memory usage (at step t:

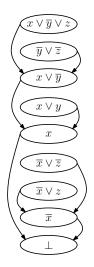
clauses before t used after t)

Width: size of the largest clause

Example:

Size 9 Space 4

Width



Space at current step

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- annotated list or
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Size: number of steps in refutation

Space: memory usage (at step t:

clauses before t used after t)

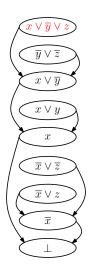
Width: size of the largest clause

Example:

Size 9

Space 4

Width 5



Space at current step

Can represent refutation as

- annotated list or
- DAG

Size: number of steps in refutation

Space: memory usage (at step t:

clauses before t used after t)

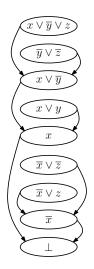
Width: size of the largest clause

Example:

Size 9

Space 4

Width 3



Space at current step

Relation Between Width and Size/Space

Width helps us understand size and space

Makes most sense for small width formulas — focus on k-CNFs

Size: Ben-Sasson and Wigderson '99

$$\log(\text{Size}) \gtrsim \text{Width}$$

Proof by syntactically manipulating short refutation into narrow refutation

Space: Atserias and Dalmau '03

More involved proof in terms of strategies for Ehrenfeucht-Fraissé games

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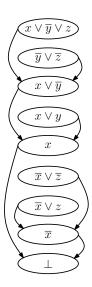
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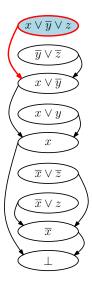
More involved proof in terms of strategies for Ehrenfeucht-Fraissé games

Our result: Simple purely syntactic proof

Razborov independently obtained a similar proof

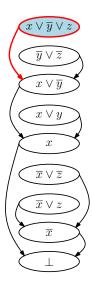






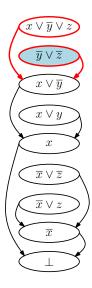


Refutation presented on whiteboard



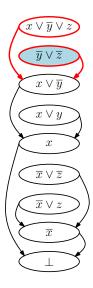


Refutation presented on whiteboard



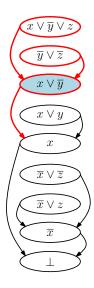
$$x \vee \overline{y} \vee z$$
$$\overline{y} \vee \overline{z}$$

Refutation presented on whiteboard



$$x \vee \overline{y} \vee z$$
$$\overline{y} \vee \overline{z}$$

Refutation presented on whiteboard

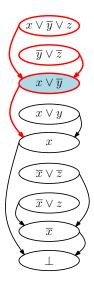


$$x \vee \overline{y} \vee z$$

$$\overline{y} \vee \overline{z}$$

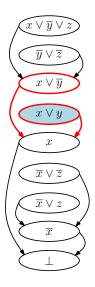
$$x \vee \overline{y}$$

- Write down axioms
- Use resolution rule



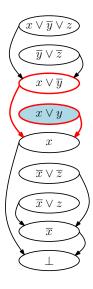
$$\begin{array}{c}
x \vee \overline{y} \vee z \\
\overline{y} \vee \overline{z} \\
x \vee \overline{y}
\end{array}$$

- Write down axioms
- Use resolution rule



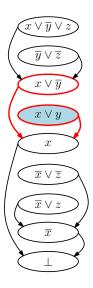
$$x \vee \overline{y} \vee z$$
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- Write down axioms
- Use resolution rule



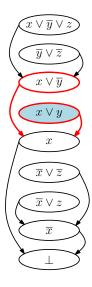


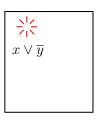
- Write down axioms
- Use resolution rule
- Erase clause



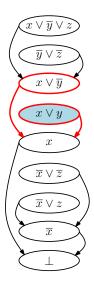


- Write down axioms
- Use resolution rule
- Erase clause



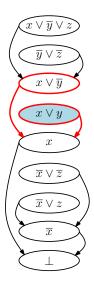


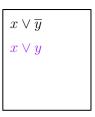
- Write down axioms
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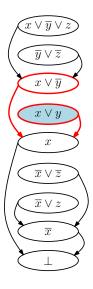


- Write down axioms
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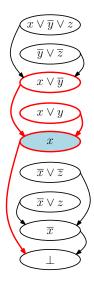
- Write down axioms
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- Write down axioms
- Use resolution rule
- Erase clause

Whiteboard Interpretation of Space

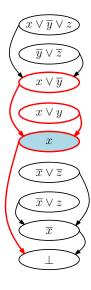




Refutation presented on whiteboard

- Write down axioms
- Use resolution rule
- Erase clause

Whiteboard Interpretation of Space



$$x \vee \overline{y}$$

$$x \vee y$$

$$x$$

Refutation presented on whiteboard

- Write down axioms
- Use resolution rule
- Erase clause

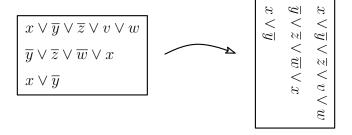
Space: max # clauses on board

Proof Idea in One Slide

No finite model theory and no Ehrenfeucht-Fraïssé games Want to turn small-height whiteboard into small-width one

Proof Idea in One Slide

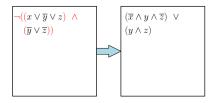
No finite model theory and no Ehrenfeucht-Fraïssé games
Want to turn small-height whiteboard into small-width one
Rotate whiteboard and get narrow whiteboard



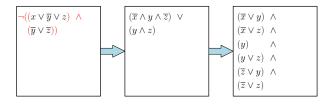
 $\begin{array}{l} x\vee\overline{y}\vee z\\ \overline{y}\vee\overline{z} \end{array}$



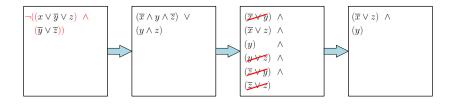
• View clauses on whiteboard as CNF and negate



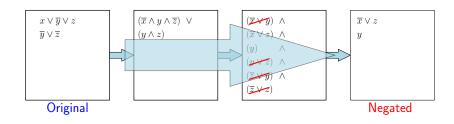
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- Apply DeMorgan's rules



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- Expand the formula into CNF by distributing OR over ANDs
- Remove trivial and redundant clauses

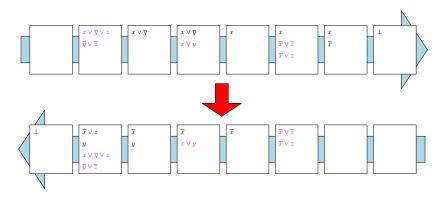


- View clauses on whiteboard as CNF and negate
- Apply DeMorgan's rules
- Expand the formula into CNF by distributing OR over ANDs
- Remove trivial and redundant clauses
- Write CNF on whiteboard

Space (# clauses) of Original \geq Width of Negated

Consequences of Negation

Negate every whiteboard and run refutation in reverse



Note: Empty whiteboard turns into contradiction and vice versa

Small space refutation is transformed into narrow one

Missing Details

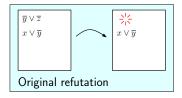
Need two things

- Prove we have backbone of resolution refutation
- Fill in missing details (without blowing up width)

Proof by case analysis over derivation steps:

- Axiom download
- Resolution rule application
- Clause erasure

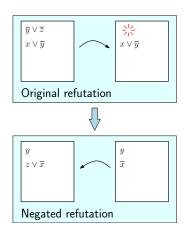
Clause Erasure



Original: Erasure weakens whiteboard

Right board weaker than left board

Clause Erasure



Original: Erasure weakens whiteboard

Right board weaker than left board

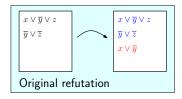
Negated: Negation inverts relation

Left board weaker than right board

Negated refutation run in reverse!

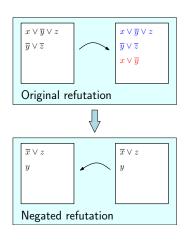
Can skip weaker whiteboards

Resolution Rule Application



Original: No change in semantic content

Resolution Rule Application

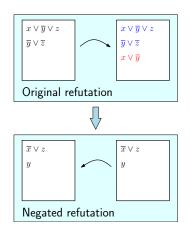


Original: No change in semantic content

Negated: No change in syntactic content (after prunning redundant clauses)

Whiteboard stays the same!

Resolution Rule Application



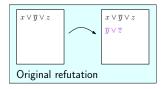
Original: No change in semantic content

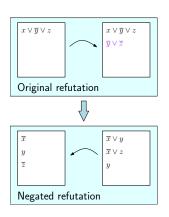
Negated: No change in syntactic content (after prunning redundant clauses)

Whiteboard stays the same!

Note: No work done thus far!

Original: Add axiom A to whiteboard

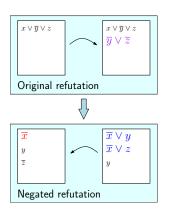




Original: Add axiom A to whiteboard

Negated: For every literal $a \in A$ add \overline{a} to all clauses of whiteboard

Use clauses $C \vee \overline{a}$ and A to derive C



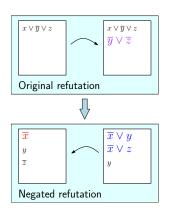
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Use clauses $C \vee \overline{a}$ and A to derive C

$$\begin{array}{c|cccc} \overline{y} \vee \overline{z} & \overline{x} \vee y \\ \hline \underline{\overline{x}} \vee \overline{z} & \overline{x} \vee z \\ \hline \overline{x} & \end{array}$$

Adds constant width to derivation



Original: Add axiom A to whiteboard

Negated: For every literal $a \in A$ add \overline{a} to all clauses of whiteboard

Use clauses $C \vee \overline{a}$ and A to derive C

$$\begin{array}{c|cccc} \overline{y} \vee \overline{z} & \overline{x} \vee y \\ \hline \underline{\overline{x} \vee \overline{z}} & \overline{x} \vee z \\ \hline \overline{x} & \end{array}$$

Adds constant width to derivation

Theorem

Space ≥ Width

Open Problem: Similar Problem for Polynomial Calculus

Polynomial calculus

Stronger proof system based on algebraic reasoning

Lines are polynomial equations instead of clauses

Degree of refutation analogous to width in resolution

Size: Impagliazzo, Pudlák, and Sgall '99

 $\log(\mathsf{Size}) \gtrsim \mathsf{Degree}$

Open Problem: Similar Problem for Polynomial Calculus

Polynomial calculus

Stronger proof system based on algebraic reasoning

Lines are polynomial equations instead of clauses

Degree of refutation analogous to width in resolution

Size: Impagliazzo, Pudlák, and Sgall '99

$$\log(\mathsf{Size}) \gtrsim \mathsf{Degree}$$

Open Problem

Is Space ≥ Degree in polynomial calculus?

Original motivation for our work

We show our approach is unlikely to work (see paper for details)

Concluding Remarks

- Space upper bounds width in resolution [Atserias and Dalmau '03]
- This work: New simple proof of this theorem
- Open problem: Space-degree relation in polynomial calculus?

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Thank you for your attention!